

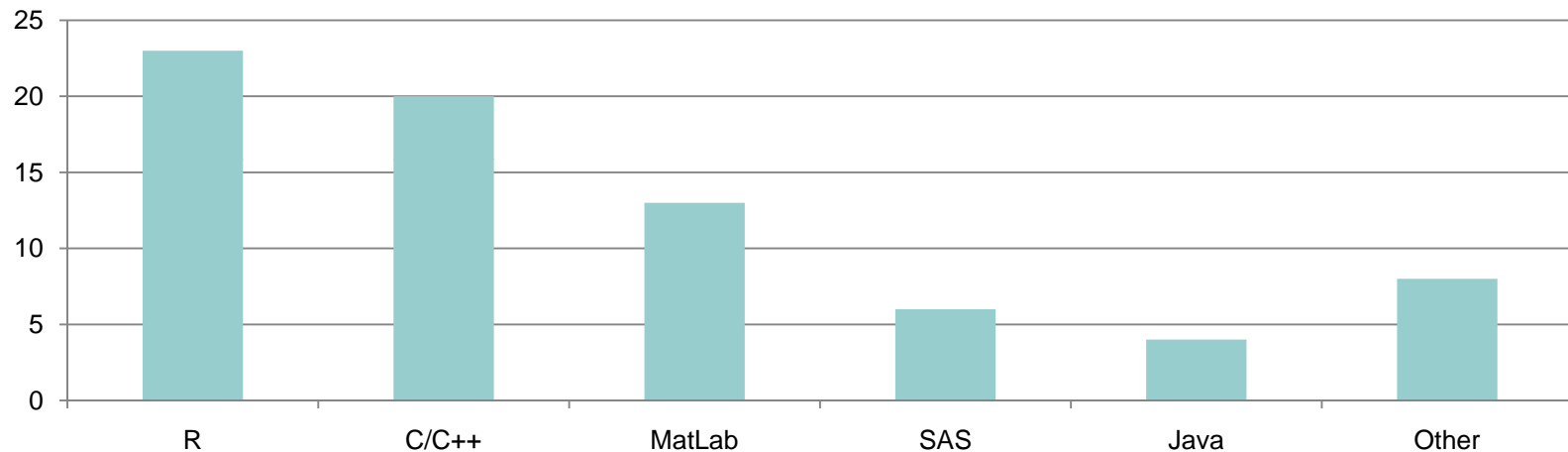
*Principles of Algorithm  
Analysis*

**Biostatistics 615/815**

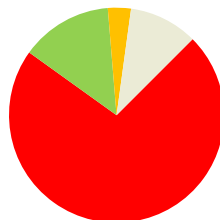
**Lecture 3**

# Snapshot of Incoming Class

## Programming Languages

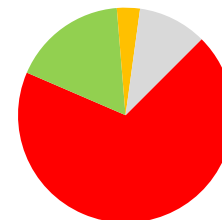


**Can you describe the QuickSort Algorithm?**



■ No  
■ Yes  
■ Maybe  
■ Blank

**Can you describe Simulated Annealing?**



■ No  
■ Yes  
■ Maybe  
■ Blank

# Homework Notes

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- Provide a hard copy including
  - Your Answers
  - Your Code
- Write specific answer to each question
  - Supported by table or graph if appropriate
- Source code
  - Indented and commented, if appropriate

# Office Hours

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- **Wednesdays**
  - 1:30 – 4:00 pm
  - SPH Tower, M4614
- **Alternatively, e-mail me at:**
  - [goncalo@umich.edu](mailto:goncalo@umich.edu)

# Last Week

## An Introduction to C

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- Strongly typed language
  - Variable and function types set explicitly
- Functional language
  - Programs are a collection of functions
- Compiling and debugging C programs
  - Setup a basic projects
  - Review compile errors and warnings
  - Step through code line by line
  - Set breakpoints

# Today

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- Strategies for comparing algorithms
- Common relationships between algorithm complexity and input data
- Compare two simple search algorithms

# Objectives

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- Framework for
  - Empirical Testing
  - Theoretical Analysis
- Highlight performance characteristics of algorithms

# Specific Questions

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- Compare two algorithms for one task
- Predict performance in a new environment
  - If we had a computer that was 10x faster and could store 10x more data, how would approach perform?
- Set values of algorithm parameters



# Two Common Mistakes

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- Ignore performance of algorithm
  - Shun faster algorithms to avoid complexity in program
  - Waiting for “simple” but inefficient algorithms to run, when efficient alternatives of modest complexity exist
- Too much weight on performance of algorithm
  - Improving program that is already very fast not worth it
  - Time spent tinkering with code is useful

# Empirical analysis

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- Given two algorithms ... which is better?
- Run both
  - Say, algorithm A takes 3 seconds
  - Say, algorithm B takes 30 seconds
- Empirical studies may not always be practical
  - Some algorithms may take too long to run!
  - Other algorithms may take too long to code...

# Choices of Input Data

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- Actual data
  - Measures performance in use
- Random data
  - Generic approach, may not be representative
- Perverse data
  - Attempt worst case analysis

# Limitations of Empirical Analysis

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- Quality of implementation
  - Is our favored implementation coded more carefully than another?
- Extraneous factors
  - Compiler
  - Machine
  - Computer system

# Limitations of Empirical Analysis

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- Requires a working program
- Theoretical analysis is an alternative
  - Estimate potential gains
- Predict effectiveness relative to new algorithms or computers (that may not yet exist)

# Theoretical Analysis

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- Predict performance of algorithm based on theoretical properties
- “Independent” of actual implementation
- Several constructs occur frequently in algorithm analysis

# Limitations of Theoretical Analysis

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- Efficiency can depend on compiler
- Efficiency may fluctuate with input data
- Some algorithms are not well understood

## The idea...

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- Given a code fragment

```
#Find parent of node i  
i = a[i];
```

- Consider how many times it is executed
- But not how long each execution takes



## Two typical analyses

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- Average-case for random input
- Worst-case
- Are these representative of real world problems?
  - Check with empirical predictions...

# The Primary Parameter $N$

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- Examples
  - Number of parameters to likelihood function
  - Number of items in dataset to be processed
  - Number of characters in a string
  - Size of file to be sorted
  - Some other abstract measure of problem size
- With multiple inputs, focus on one at a time, while holding the others constant

## Running time as a function of $N$

$f(N)$	Description	Running time when $N$ doubles...
$1$	<i>constant</i>	-
$\log N$	<i>logarithmic</i>	constant increase
$N$	<i>linear</i>	doubles
$N \log N$	<i>log-linear</i>	more than doubles
$N^2$	<i>quadratic</i>	increases fourfold
$N^3$	<i>cubic</i>	increases eightfold
$2^N$	<i>exponential</i>	running time squares

# Running time as a function of $N$

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- Multiple terms may be involved
  - e.g.  $N + N \log N$
- Typically, we ignore
  - Smaller terms
  - Constant coefficient
  - Focus on inner loop
- In rare cases, smaller terms and constant coefficient will be important

# Time to Solve Large Problem

operations per second	Problem Size $N = 1,000,000$		
	$N$	$N \log N$	$N^2$
$10^6$	seconds	minutes	months
$10^9$	instant	instant	hours
$10^{12}$	instant	instant	seconds

# Time to Solve Huge Problem

operations per second	Problem Size $N = 1,000,000,000$		
	$N$	$N \log N$	$N^2$
$10^6$	hours	days	never
$10^9$	seconds	minutes	centuries
$10^{12}$	instant	instant	months

# Application

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- Analysis of two search algorithms
- Each algorithm:
  - Considers a set of items stored in an array
  - Searches through items to decide whether a particular value occurs

# Sequential Search

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```
int search(int a[], int value, int start, int stop)
{
    // Variable declarations
    int i;

    // Search through each item
    for (i = start; i <= stop; i++)
        if (value == a[i])
            return i;

    // Search failed
    return -1;
}
```



# Sequential Search Properties

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- Algorithm:
  - Look through array sequentially, until we find a match
- Average cost
  - If match found:  $N/2$
  - If match not found:  $N$
- Actual cost depends on fraction of successful searches

## Better Sequential Search

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- If items are sorted...
- Stop unsuccessful search early, when we reach item with higher value
  - Cost for unsuccessful searches is now  $N/2$
- Overall, algorithm is still  $O(N)$

# Binary Search

```
int search(int a[], int value, int start, int stop)
{
    while (stop >= start)
    {
        // Find midpoint
        int mid = (start + stop) / 2;

        // Compare midpoint to value
        if (value == a[mid])
            return mid;

        // Reduce input in half !...
        if (value > a[mid])
            { start = mid + 1; }
        else
            { stop = mid - 1; }
    }

    // Search failed
    return -1;
}
```

# Binary Search Properties

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- Algorithm:
  - Halve number of items to consider with each comparison
- Worst-case cost
  - Maximum cost is never greater than  $\log_2 N$
- Much better than sequential search, but even better methods exist!

## Sequential vs. Binary Search

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N	<u>M = 1,000</u>		<u>M = 10,000</u>		<u>M = 100,000</u>	
	S	B	S	B	S	B
125	1	1	13	2	130	20
250	3	0	25	2	251	22
500	5	0	49	3	492	23
1250	13	0	128	3	1276	25
2500	26	1	267	3	*	28

Timings in seconds, for M searches in table of N elements

# Big-Oh Notation

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- Algorithm is  $O(N)$  or  $O(N \log N)$ 
  - Common statement
  - What does it mean?
- Summarizes performance for large  $N$
- Focuses on leading terms of expression describing running time

## Big-Oh Notation

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- Consider function  $g(N)$
- It is said to be  $O(f(N))$
- If there exist  $c_0$  and  $N_0$  such that:
  - $N > N_0$  implies  $c_0 f(N) > g(N)$

## From N to Running Time...

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- Common relationships
  - $N^2$
  - $\log N$
  - $N \log N$
  - $N$
- Describe examples of how these arise
- Cost of running program is  $C_N$



$O(N^2)$

- Loop through input successively, eliminate one item at a time

$$\begin{aligned}C_N &= C_{N-1} + N \quad \text{for } N \geq 2, C_1 = 1 \\ &= C_{N-2} + (N-1) + N \\ &\dots \\ &= 1 + 2 + \dots + (N-1) + N \\ &= \frac{N(N+1)}{2}\end{aligned}$$

$O(\log N)$

- Recursive program, halves input in one step

$$C_{2^n} = C_{2^{n-1}} + 1 \quad \text{for } N \geq 2, C_1 = 1$$

$$= C_{2^{n-2}} + 1 + 1$$

$$= C_{2^{n-3}} + 3$$

...

$$= C_{2^0} + n$$

$$= n + 1$$

$$N = 2^n$$

## $O(N \log N)$

- Recursive program, processes each item, splits input into two halves, examines each one...

$$C_N = 2C_{N/2} + N \quad \text{for } N \geq 2, C_1 = 0$$

$$C_{2^n} = 2C_{2^{n-1}} + 2^n$$

$$\frac{C_{2^n}}{2^n} = \frac{2C_{2^{n-1}} + 2^n}{2^n}$$

$$= \frac{C_{2^{n-1}}}{2^{n-1}} + 1$$

$$= \frac{C_{2^{n-2}}}{2^{n-2}} + 1 + 1$$

...

$$= n$$

$O(2N)$

- Halves input, must examine each item...

$$C_N = C_{N/2} + N \quad \text{for } N \geq 2, C_1=1$$

$$= N + \frac{N}{2} + \frac{N}{4} + \frac{N}{8} + \dots$$

$$\approx 2N$$

## Summary

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- Outline principles for analysis of algorithms
- Introduced some common relationships between  $N$  and running time
- Described two simple search algorithms

## Further Reading

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- Read chapter 2 of Sedgewick

## Tip of the Day: Defensive Programming

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- Document code and programs
  - Indicate intended purpose
  - Specify required inputs
  - Always indicate author
- Check for error conditions