Course Objective

- Provide students with a grounding for conducting statistical computing projects
- Applications and examples will be in C and in R.
- But the focus is on an algorithmic way of thinking!
Part I: Key Algorithms

- Connectivity
- Sorting
- Searching
- Hashing
- Key data structures
Part II: Statistical Methods

- Random Numbers
- Markov-Chain Monte-Carlo
  - Metropolis-Hastings
  - Gibbs Sampling
- Function Optimization
  - Naïve algorithms
  - Newton’s Methods
  - E-M algorithm
- Numerical Integration
Textbooks

- Algorithms in C

- Numerical Recipes in C
Assessment for 615

- Weekly Assignments
  - About 60% of the final mark

- 2 Exams
  - About 40% of the final mark
Assessment for 815

- Weekly Assignments
  - About 40% of the final mark

- 2 Exams
  - About 30% of the final mark

- Group Project
  - About 30% of the final mark
Office Hours

- Wednesdays
  - 10:00 – 11:00 am

- Room M4132
  School of Public Health II
Algorithms

- Methods for solving problems that are well suited to computer implementation
- Good algorithms make apparently impossible problems become simple
Algorithms are ideas ...

- Focus on approach to a problem

- Typically, the actual implementation could be take many different forms
  - Computer languages
  - Pen and paper
Example: DNA Sequence Matches

- When the Human Genome Project started, searching through the entire genome sequence seemed impractical...

- For example,
  - Searching for ~150 sequences of about 500bp each in ~3,000,000,000 bases of sequence would take ~3 hours with the BLAST or FASTA3 algorithms
Example: DNA Sequence Matches

- Mullikin and colleagues (2001) described an improved algorithm, using hash tables, that could do this in < 2 seconds

- Reference:
Today’s Lecture

- Introduce a “Connectivity problem” and some alternative solutions

- If you haven’t done much programming before, don’t worry too much about implementation details.
  - We’ll fill these in later lectures.
The Connectivity Problem

- N objects
  - Integer names 0 .. N – 1

- M connections between pairs of objects
  - Each connection identifies a pair \((p, q)\)

Possible questions:
- Are all objects connected?
- Are some connections redundant?
Possible applications

- Is a direct connection between two computers required in a network?
  - Or can we use some existing connections instead?

- Are two individuals part of the same extended family in a genetic study?
Are the two points connected?
Specific Question

- Can we identify redundant connections?
  - A redundant connection would link two points that are already connected

- For N objects there can be no more than N-1 non-redundant connections
  - Corresponds to all points being connected
A simple example ...

- Connections
  - 3-4
  - 4-9
  - 8-0
  - 2-3
  - 5-6
  - 2-9
  - 4-8
  - 0-2
A simple example ...

- Connections
  - 3-4 ✓
  - 4-9 ✓
  - 8-0 ✓
  - 2-3 ✓
  - 5-6 ✓
  - 2-9 Redundant: 2-3 ; 3-4 ; 4-9
  - 4-8 ✓
  - 0-2 Redundant: 0-8; 8-4; 4-3; 3-2
Specific Tasks

As we proceed through list of connections, conduct two tasks:

- Decide if each connection is new.
- Incorporate information about new connections.
The Fundamental Operations

- The *Find* operation
  - Identify the set containing a particular object or items.

- The *Union* operation
  - Replace the sets containing two groups of objects by their union.
The First Step

- Developing a solution that works correctly
  - Should be simple
  - Easy to verify correctness
  - May not be most efficient

- Useful as check of “better” solutions…
Arrays of Integers

- Simple data structure
  - Analogous to a vector

- The notation $a[i]$ refers to the $i^{th}$ integer in the array
  - We’ll typically pre-specify the total number of integers
Quick Find Algorithm

- **Data**
  - Array of N integers
  - Objects \( p \) and \( q \) connected iif \( a[p] == a[q] \)

- **Setup**
  - Initialize \( a[i] = i \), for \( 0 \leq i < N \)

- **For each pair**
  - If \( a[p] == a[q] \) objects are connected (FIND)
  - Move all entries in set \( a[p] \) to set \( a[q] \) (UNION)
A Simple C Implementation

#define N 1000

int main()
{
    int i, p, q, set, a[N];  // Variable declarations

    for (i = 0; i < N; i++)  // Data Initialization
        a[i] = i;

    while (scanf(" %d %d", &p, &q) == 2) // Loop through connections
    {
        if (a[p] == a[q]) continue;  // FIND operation

        set = a[p];  // UNION operation
        for (i = 0; i < N; i++)
            if (a[i] == set)
                a[i] = a[q];

        printf("%d %d is a new connection\n", p, q);
    }
    return 0;
}
Pictorial Representation

- Array as connections are added:
  - 3-4
  - 4-9
  - 8-0
  - 2-3
  - 2-9 * Redundant *

![Diagram showing connections between numbers 0 to 9]

- Connections added:
  - 3-4
  - 4-9
  - 8-0
  - 2-3
  - 2-9 (marked as redundant)
How efficient is Quick Find?

- If there $N$ objects and $M$ connections*, the Quick Find algorithm requires on the order of $MN$ operations.

- Not feasible for very large numbers of objects...

* In this case only non-redundant connections actually count.
Quick-Union Algorithm I

- Complementary to Quick Find
- More complex data organization
  - Each object points to “parent” object in the same set
Quick-Union Algorithm II

- For each pair
  - Follow pointers until we reach object that points to itself
  - If $a[p]$ and $a[q]$ eventually lead to the same object, we are in the same set (FIND)
  - Otherwise, link the object to which $a[p]$ leads to the object which $a[q]$ leads (UNION)
C Implementation

// Loop through connections
while (scanf(" %d %d", &p, &q) == 2) {
    // Check that input is within bounds
    if (p < 0 || p >= N || q < 0 || q >= N) continue;

    // FIND operation
    for (i = a[p]; a[i] != i; i = a[i]) ;
    for (j = a[q]; a[j] != j; j = a[j]) ;
    if (i == j) continue;

    // UNION operation
    a[i] = j;

    printf("%d %d is a new connection\n", p, q);
}
Pictorial Representation

- Array as connections are added:
  - 3-4
  - 4-9
  - 8-0
  - 2-3
  - 2-9 * Redundant
How efficient is Quick Union?

- Quick Union is typically faster than Quick Find.

- However, the data can conspire to make things difficult:
  - If objects were paired 1-2; 2-3; 3-4; 4-5; … we’ll build long chains which slow down FIND operations

- In the worst case, we can still need about MN operations
Weighted Quick Union

- A smarter version of Quick Union, that avoids long chains

- Keep track of the number of elements in each set (using a separate array)

- Link smaller set to larger set
  - Union increases length of chains in smaller set by 1
C Implementation

// Initialize weights
for (i = 0; i < N; i++)
    weight[i] = i;

// Loop through connections
while (scanf(" %d %d", &p, &q) == 2)
{
    // Check that input is within bounds
    if (p < 0 || p >= N || q < 0 || q >= N) continue;

    // FIND operation
    for (i = a[p]; a[i] != i; i = a[i]) ;
    for (j = a[q]; a[j] != j; j = a[j]) ;
    if (i == j) continue;

    // UNION operation
    if (weight[i] < weight[j])
        { a[i] = j; weight[j] += weight[i]; }
    else
        { a[j] = i; weight[i] += weight[j]; }

    printf("%d %d is a new connection\n", p, q);
}
Pictorial Representation

- Array as connections are added:
  - 3-4
  - 4-9
  - 8-0
  - 2-3
  - 2-9 * Redundant
Efficiency of Weighted Quick Union

- Guarantees that pointer chains are no more than $\log_2 N$ elements long
- Overall, requires about $M \log_2 N$ operations
- Suitable for very large data sets with millions of objects and connections
Pictorial Comparison
Quick Union  Quick Find  Weighted
<table>
<thead>
<tr>
<th>Nodes (Connections)</th>
<th>Quick Find</th>
<th>Quick Union</th>
<th>Weighted Quick Union</th>
</tr>
</thead>
<tbody>
<tr>
<td>50,000 (50,000)</td>
<td>6</td>
<td>1</td>
<td>&lt;1</td>
</tr>
<tr>
<td>100,000 (100,000)</td>
<td>12</td>
<td>4</td>
<td>&lt;1</td>
</tr>
<tr>
<td>200,000 (200,000)</td>
<td>25</td>
<td>15</td>
<td>&lt;1</td>
</tr>
</tbody>
</table>
Summary

- Considered 3 alternative solutions to the “connectivity problem”
  - Are any connections in a set redundant?
  - Are all objects in a set connected?

- Compared some of the computational cost for the different methods
Reading Material

- Read Chapter 1 of Sedgewick